

The Card Guessing Game: A generating function approach

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Introduction

Blackjack: Riffle shuffle. As a player, we do card counting.

In this project (similar to Blackjack):

Original deck is $[1, 2, 3, \dots, n]$.

Riffle shuffle k times, guess card \rightarrow reveal card, then repeat.

We concentrate on the optimal guessing strategy and statistics on the number of correct guesses.

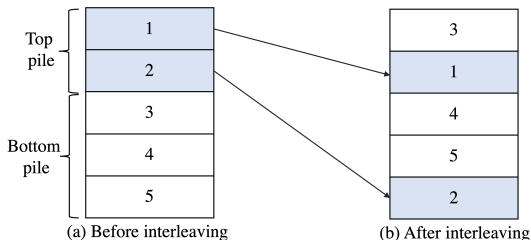
Gilbert-Shannon-Reeds (GSR) Model for Riffle Shuffles

1. Split the deck into two piles.

The probability of cutting the top t cards is $\frac{\binom{n}{t}}{2^n}$.

2. Then, interleave the piles back into a single one.

Each interleaving has probability $\frac{1}{2^n}$ to come up.



Example of 1-time riffle shuffle of a deck of 5 cards

1-time riffle shuffle: examples

$n=1$ 1 1

$n=2$ 1 1 1 2
 2 2 2 1

$n=3$ 1 1 1 1 1 2 2 3
 2 2 2 2 3 1 3 1
 3 3 3 3 2 3 1 2

$n=4$ 1 1 1 1 1 1 1 1
 2 2 2 2 2 2 3 3
 3 3 3 3 3 4 2 4
 4 4 4 4 4 3 4 2

 1 2 2 2 3 3 3 4
 4 1 3 3 1 1 4 1
 2 3 1 4 2 4 1 2
 3 4 4 1 4 2 2 3

Optimal guessing strategy

Algorithm:

- 1: Start by guessing number 1.
- 2: If true then continue to guess the next number in line.
- 3: If false then the deck is now split into two increasing subsequences. Guess the first element in the longer subsequence.
- 4: Continue to guess this way until until no cards remain.

This algorithm is proved to provide the maximum expected number of correct guesses.

Example

Permutation π	1	1	1	1	1	2	2	2	3	3	3	4
	2	2	3	3	4	1	3	3	1	4	1	1
	3	4	2	4	2	3	1	4	2	1	4	2
	4	3	4	2	3	4	4	1	4	2	2	3
$P(\pi)$	5/16	1/16	1/16	1/16	1/16	1/16	1/16	1/16	1/16	1/16	1/16	1/16
#Correct guesses	4	3	3	2	3	2	3	2	3	2	2	3

All possible permutations after shuffling a 4-card deck once. The color indicates a correct guess under the optimal strategy.

Goal

The goal is to calculate the moments (i.e. mean, variance, etc) of the number of correct guesses (denote X_n) amongst all of the resulting permutations.

Generating functions and recurrences

Generating functions:

$$D_n(q) = \sum_{i=0}^{\infty} a_i q^i,$$

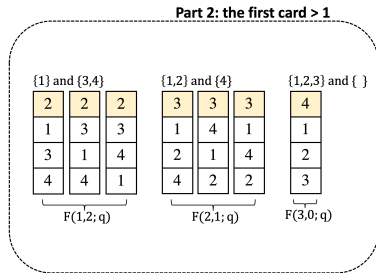
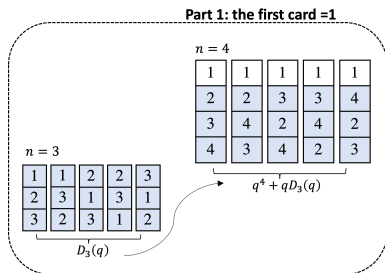
where a_i denotes the number of permutations with i correct guesses.

Recurrence:

$$D_n(q) = \underbrace{qD_{n-1}(q) + q^n}_{\text{the first card} = 1} + \underbrace{\sum_{i=0}^{n-2} F(n-1-i, i; q)}_{\text{the first card} > 1}, \quad (\text{Main recurrence})$$

where $D_0(q) = 1$.

Example:



$$\text{Recurrence structure } D_4(q) = (q^4 + qD_3(q)) + F(1,2; q) + F(2,1; q) + F(3,0; q)$$

The catch!

$$F(m, n; q) = \underbrace{qF(m-1, n; q)}_{\text{next card from longer subseq.}} + \underbrace{F(m, n-1; q)}_{\text{next card from shorter subseq.}}, \quad (1)$$

for $m \geq n$, where $F(m, 0; q) = q^m$. Also, $F(m, n; q) := F(n, m; q)$ whenever $m < n$.

It is easy to show the formula of $F(m, n; q)$ (once you know what it looks like).

Proposition

For $m \geq n$,

$$F(m, n; q) = \sum_{i=0}^n \left[\binom{m+n}{i} - \binom{m+n}{i-1} \right] q^{m+n-i}. \quad (2)$$

More catch!!

Assume $G_n(q) = q^n + \sum_{i=0}^{n-2} F(n-1-i, i; q)$.

It can be shown that

Proposition

For $r \geq 1$, the closed-form formula for $G_n^{(r)}(q)|_{q=1}$ can be obtained by evaluating the binomial sums:

$$G_{2k}^{(r)}(q)|_{q=1} = (2k)_r - (2k-1)_r + 2 \sum_{i=0}^{k-1} (k-i) \left[\binom{2k-1}{i} - \binom{2k-1}{i-1} \right] (2k-1-i)_r,$$

$$G_{2k+1}^{(r)}(q)|_{q=1} = (2k+1)_r - (2k)_r + 2 \sum_{i=0}^k (k + \frac{1}{2} - i) \left[\binom{2k}{i} - \binom{2k}{i-1} \right] (2k-i)_r,$$

where $(a)_r$ is the falling factorial, i.e. $(a)_r = a(a-1)(a-2)\dots(a-r+1)$.

Now the moments!

Procedure: Factorial Moment (fixed r , formula in n)

Step 1: Compute $G_n^{(r)}(q)|_{q=1}$ by the binomial sum, n symbolic.

Step 2: Use the method of undetermined coefficient to calculate $D_n^{(r)}(q)|_{q=1}$.

Step 3: Apply (4) to obtain $E[X(X-1)\dots(X-r+1)]$.

- Step 2 is acquired through the relation:

$$D_n(q) = qD_{n-1}(q) + G_n(q). \quad (3)$$

- Step 3 is acquired from the relation:

$$E[X(X-1)\dots(X-r+1)] = \frac{D_n^{(r)}(q)|_{q=1}}{2^n}. \quad (4)$$

Some results:

For $n = 2k$:

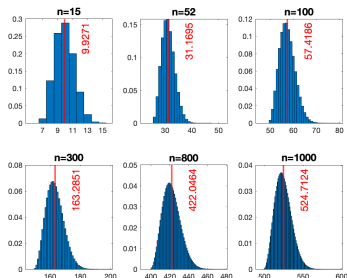
$$E[X] = \frac{n}{2} + \sqrt{\frac{2n}{\pi}} - \frac{1}{2} - \sqrt{\frac{2}{\pi n}} \left(\frac{3}{4} + \frac{49}{96n} + \frac{439}{384n^2} + \frac{76709}{18432n^3} + \dots \right).$$

$$E[(X - \mu)^2] = \left(\frac{3}{4} - \frac{2}{\pi} \right) n - \frac{3}{4} + \frac{3}{\pi} - \sqrt{\frac{2}{\pi n}} + \frac{11}{12\pi n} + \dots$$

$$E[(X - \mu)^3] = \sqrt{\frac{2}{\pi}} \left(\left(\frac{4}{\pi} - \frac{5}{4} \right) n^{3/2} + \left(\frac{43}{16} - \frac{9}{\pi} \right) n^{1/2} - \frac{3\sqrt{2\pi}}{4} + 3\sqrt{\frac{2}{\pi}} + \dots \right).$$

Non-normal distribution

The skewness coefficient is given by $\frac{m_3}{m_2^{3/2}}$, where $m_r := E[(X - \mu)^r]$. We see that the skewness of X_n does not tend to zero. Therefore, the number of correct guesses is not asymptotically normally distributed.







Probability histograms of the number of correct guesses when n varies. The red vertical line indicates the corresponding expected value $E[X_n]$.




Generalization to k Riffle Shuffles: Is it possible?

This problem seems very difficult at the moment.

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