

# Statistics of Domino Tilings on a Rectangular Board

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# Introduction

It is well known the Fibonacci sequence,  $F_n$ , is the number of ways to cover a 2-by- $(n - 1)$  board using only the horizontal( $H$ ) or vertical( $V$ ) 2-by-1 dominos.

It is natural to generalize this idea to a rectangular  $m$ -by- $n$  board where  $m$  is a fixed number and  $n$  is symbolic.

We can try harder and consider the mixed moment  $E[V^a H^b]$  for fixed non-negative integers  $a, b$  but general  $m, n$ . After all these moments will give an information of the distribution of “ $V$ - $H$  statistics”.

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# Fibonacci numbers and Statistics on an 2-by- $n$ board

We fixed the board to size 2-by- $n$ .

Let  $V_n$  be a random variable of the number of vertical dominos on the tiling of the board of size 2-by- $n$ .

First consider the power sum,  $S[V^r] := \sum_{b \in B_n} V(b)^r$ .

For  $r = 0$ , the number of possible tilings on the 2-by- $n$  board,  $S[V^0] = f_n$ , where  $f_n$  is the  $(n + 1)$ -th Fibonacci numbers.

Next consider the straight moment,

$$E[V^r] = S[V^r]/S[V^0] = S[V^r]/f_n.$$

The average number of  $V$ , vertical domino, on a 2-by- $n$  board,

$$\mu_n := E[V] = S[V]/f_n.$$

# Fibonacci numbers and Statistics on an $2$ -by- $n$ board

Eventually we are interested in the moment about the mean,  $E[(V - \mu)^r]$  and then the scaled-moments  $\frac{E[(V - \mu)^r]}{E[(V - \mu)^2]^{r/2}}$  (which can show the normality of an asymptotic distribution).

## Fast Calculation

The generating function of a random variable  $V$  is defined by

$$F_n(v) := \sum_{b \in B_n} v^{V(b)}.$$

For example, on the board of size 2-by-3,  $F_3(v) = v^3 + 2v$ .

$S[V^0] = F_n(1)$  and  $S[V^r]$  is obtained by applying the operator  $(v \frac{d}{dv})^r$  to  $F_n(v)$ , then substitute  $v = 1$ .

## Fast Calculation

Next, we define the grand generating function  $H(v, t)$  by

$$H(v, t) = \sum_{n=0}^{\infty} F_n(v) t^n.$$

In this problem, we have

$$H(v, t) = \frac{1}{1 - vt - t^2}, \quad (1)$$

which can be derived from a simple recurrence

$$H(v, t) = 1 + vtH(v, t) + t^2H(v, t).$$



# Fast Calculation

With equation (1) in hand, we calculate data of  $S[V^r]$  (and hence  $E[V^r]$ ) very fast. Recall that

$$\sum_{n=0}^{\infty} S[V^0]t^n = H(1, t) = \frac{1}{1 - t - t^2}.$$

In general,

$$\sum_{n=0}^{\infty} S[V^r]t^n = \left(v \frac{d}{dv}\right)^r H(v, t)|_{v=1}.$$

## Fast Calculation

From the quotient rule in calculus, we find the fact that

$$\sum_{n=0}^{\infty} S[V^r] t^n = \left( v \frac{d}{dv} \right)^r H(v, t) \Big|_{v=1} = \frac{P_r(t)}{(1-t-t^2)^{r+1}}.$$

where  $P_r(t)$  is a polynomial in  $t$  of degree at most  $2r$ .

Hence  $S[V^r]$  in fact satisfies the recurrence of the form  $(N^2 - N - 1)^{r+1}$  and can be written in the form,

$$S[V^r] = A(n)f_n + B(n)f_{n-1}$$

where  $A(n)$  and  $B(n)$  are polynomial in  $n$  of degree at most  $r$ .

## Conjectures

We then be able to use computer program to conjecture these formulas by trying to fit the polynomial to the data. Given  $\phi = \frac{1 + \sqrt{5}}{2}$ .

$$f_n := S[V^0] = \frac{1}{\sqrt{5}} \left( \frac{1 + \sqrt{5}}{2} \right)^{n+1} - \frac{1}{\sqrt{5}} \left( \frac{1 - \sqrt{5}}{2} \right)^{n+1}.$$

$$\mu_n := E[V] = \frac{1}{f_n} \left[ \frac{n}{5} f_n + \frac{2}{5} (n+1) f_{n-1} \right] \approx \frac{n}{5} + \frac{2}{5\phi} (n+1),$$

$$\begin{aligned} E[V^2] &= \frac{1}{f_n} \left[ \frac{n(5n+12)}{25} f_n + \frac{4(n+1)}{25} f_{n-1} \right] \\ &\approx \frac{n(5n+12)}{25} + \frac{4(n+1)}{25\phi}. \end{aligned}$$

## Conjectures

From these calculations, we are able to make a general conjecture of  $E[V^r]$  starting from the leading terms (and after simplify  $\phi$ ):

### Conjectures

$E[V^r]$

$$\approx \frac{n^r}{5^{r/2}} \left[ 1 + \frac{r(2r + \sqrt{5} - 3)}{\sqrt{5}n} + \frac{r(r-1)(6r^2 - 32r + 6\sqrt{5}r + 37 - 9\sqrt{5})}{15n^2} + \frac{2\sqrt{5}r(r-1)(r-2)(2r^3 + 3\sqrt{5}r^2 - 23r^2 - 16\sqrt{5}r + 77r + 16\sqrt{5} - 73)}{75n^3} + \frac{r(r-1)(r-2)(r-3)(6r^4 - 120r^3 + 12\sqrt{5}r^3 - 138\sqrt{5}r^2 + 800r^2 - 2100r + 225n^4)}{225n^4} + \dots \right],$$

## Conjectures

We use the conjectures of straight moments to calculate the moment about the mean:

$$E[(V - \mu)^0] = 1,$$

$$E[(V - \mu)^1] = 0,$$

$$E[(V - \mu)^2] = E[X^2] - E[X]^2 = \frac{4\sqrt{5}n + 4\sqrt{5} - 8}{25},$$

$$E[(V - \mu)^3] = E[X^3] - 3E[X^2]E[X] + 2E[X]^3 = \frac{8\sqrt{5}n}{125} + \frac{8\sqrt{5}}{125} - \frac{48}{125},$$

$$E[(V - \mu)^4] = \frac{48n^2}{125} + \frac{96n}{125} - \frac{272\sqrt{5}n}{625} + \frac{16}{25} - \frac{272\sqrt{5}}{625},$$

$$E[(V - \mu)^5] \approx \frac{64n^2}{125} + \frac{128n}{125} - \frac{736\sqrt{5}n}{625} + \frac{9776}{3125} - \frac{144\sqrt{5}}{125},$$

$$E[(V - \mu)^6] \approx \frac{192\sqrt{5}n^3}{625} + \dots,$$

# Conjectures

These data lead us to the conjectures of general formulas of moment about the mean which we will prove formally.

$$E[(V - \mu)^{2r}] = \left(\frac{2}{5\sqrt{5}}\right)^r \frac{(2r)!n^r}{r!} + \text{smaller terms,}$$

$$E[(V - \mu)^{2r+1}] = \frac{2}{15} \left(\frac{2}{5\sqrt{5}}\right)^r \frac{(2r+1)!n^r}{(r-1)!} + \text{smaller terms.}$$

## Toward the Proof

We will prove the conjectures on  $E[V^r]$  and  $E[(V - \mu)^r]$  from the previous slide. Then use them to conclude the asymptotic normality distribution of  $V$ . The method of guess and check will again play an important role here.

The generating function  $F_n(v)$  can be defined by the recurrence

$$F_n(v) = vF_{n-1}(v) + F_{n-2}(v) \quad , n \geq 2,$$

where

$$F_0(v) = 1, \quad F_1(v) = v.$$

## Toward the Proof

The centralized probability generating function of  $F_n(v)$  is

$$G_n(v) := \sum_i p(i)v^{i-\mu} = \frac{1}{f_n v^\mu} F_n(v),$$

The recurrence will look like

$$G_n(v) = v \frac{f_{n-1} G_{n-1}(v)}{f_n v^{\mu_n - \mu_{n-1}}} + \frac{f_{n-2} G_{n-2}(v)}{f_n v^{\mu_n - \mu_{n-2}}}, \quad n \geq 2, \quad (2)$$

where

$$G_0(v) = 1, \quad G_1(v) = 1.$$

We will use this recurrence to set up some relations for the proof.



## Probability Background Interim

But before diving into the proof, we need to discuss some necessary probability background.

Definition (Exponential moment generating function)

$$\phi(t) := E[e^{tX}] = \begin{cases} \sum_x e^{tx} p(x) & \text{if } x \text{ is discrete} \\ \int_{-\infty}^{\infty} e^{tx} f(x) dx & \text{if } x \text{ is continuous.} \end{cases}$$

Note  $m_n := E[X^n] = \phi^n(0)$ .

# Probability Background Interim

We then have that

$$\phi(t) := E[e^{tX}] = \sum_x \sum_n \frac{t^n x^n}{n!} p(x) = \sum_n \frac{t^n}{n!} \sum_x x^n p(x) = \sum_n \frac{t^n}{n!} m_n.$$

Moment of the standard normal distribution:

$$\phi(t) = e^{\frac{t^2}{2}} = \sum_r \frac{t^{2r}}{r!2^r} \text{ implies that } m_{2r} = \frac{(2r)!}{r!2^r} \text{ and } m_{2r-1} = 0.$$

# Main Theorem

Back to the main theorem

## Theorem

Let  $E_r(n) = \frac{E[(V - \mu)^r]}{r!}$  then

$$E_{2r}(n) = \left(\frac{2}{5\sqrt{5}}\right)^r \frac{n^r}{r!} + \text{smaller terms},$$

$$E_{2r+1}(n) = \frac{2}{15} \left(\frac{2}{5\sqrt{5}}\right)^r \frac{n^r}{(r-1)!} + \text{smaller terms}$$

# The Proof

**Proof** We see that

$$G_n(e^t) = \sum_i p(i)e^{t(i-\mu)} = \phi(t), \quad \text{where the random variable } X = i - \mu.$$

Now define the Maclaurin series of  $G_n(e^t)$  by

$$G_n(e^t) = \sum_r E_r(n)t^r.$$

By the fact of probability generating function mentioned earlier,

$$E_r(n) = \frac{E[(V - \mu)^r]}{r!}.$$

# The Proof

The recurrence (2) becomes:

$$G_n(e^t) = e^t \frac{f_{n-1} G_{n-1}(e^t)}{f_n e^{t(\mu_n - \mu_{n-1})}} + \frac{f_{n-2} G_{n-2}(e^t)}{f_n e^{t(\mu_n - \mu_{n-2})}}, \quad n \geq 2. \quad (3)$$

We apply an induction on  $r$  using (3) to show the assertion of  $E_r(n)$ .

$$\begin{aligned} G_n(e^t) &= \frac{f_{n-1}}{f_n} \left[ \sum_{r=0}^{\infty} \frac{(1 - \mu_n + \mu_{n-1})^r t^r}{r!} \right] \cdot \left[ \sum_{r=0}^{\infty} E_r(n-1) t^r \right] \\ &+ \frac{f_{n-2}}{f_n} \left[ \sum_{r=0}^{\infty} \frac{(-\mu_n + \mu_{n-2})^r t^r}{r!} \right] \cdot \left[ \sum_{r=0}^{\infty} E_r(n-2) t^r \right]. \end{aligned}$$

# The Proof

By comparing coefficient of  $t^r$ , we obtain the relations:

$$\begin{aligned}
 E_r(n) &- \frac{f_{n-1}}{f_n} E_r(n-1) - \frac{f_{n-2}}{f_n} E_r(n-2) \\
 &= \frac{f_{n-1}}{f_n} a_n E_{r-1}(n-1) + \frac{f_{n-2}}{f_n} b_n E_{r-1}(n-2) \\
 &+ \frac{f_{n-1}}{f_n} \frac{a_n^2}{2} E_{r-2}(n-1) + \frac{f_{n-2}}{f_n} \frac{b_n^2}{2} E_{r-2}(n-2) \\
 &+ \frac{f_{n-1}}{f_n} \frac{a_n^3}{6} E_{r-3}(n-1) + \frac{f_{n-2}}{f_n} \frac{b_n^3}{6} E_{r-3}(n-2) + \dots
 \end{aligned}$$

where  $a_n = 1 - \mu_n + \mu_{n-1} \approx 1 - \frac{1}{\sqrt{5}}$  and  $b_n = -\mu_n + \mu_{n-2} \approx -\frac{2}{\sqrt{5}}$ .

# The Proof

## Case1: even case:

Left hand side:

$$\begin{aligned}
 & E_{2r}(n) - \frac{f_{n-1}}{f_n} E_{2r}(n-1) - \frac{f_{n-2}}{f_n} E_{2r}(n-2) \\
 &= \left( \frac{f_{n-1}}{f_n} + 2 \frac{f_{n-2}}{f_n} \right) \left( \frac{2}{5\sqrt{5}} \right)^r \frac{n^{r-1}}{(r-1)!} + \text{smaller terms}
 \end{aligned}$$

# The Proof

Right hand side is

$$\begin{aligned} & \frac{f_{n-1}}{f_n} a_n E_{2r-1}(n-1) + \frac{f_{n-2}}{f_n} b_n E_{2r-1}(n-2) + \\ & \frac{f_{n-1}}{f_n} \frac{a_n^2}{2} E_{2r-2}(n-1) + \frac{f_{n-2}}{f_n} \frac{b_n^2}{2} E_{2r-2}(n-2) + \text{smaller terms} \\ & = \left( \frac{f_{n-1}}{f_n} \frac{a_n^2}{2} + \frac{f_{n-2}}{f_n} \frac{b_n^2}{2} \right) \left( \frac{2}{5\sqrt{5}} \right)^{r-1} \frac{n^{r-1}}{(r-1)!} + \text{smaller terms.} \end{aligned}$$

which are equal.



# The Proof

**Case2: odd case:** ... (can be done similarly).

# Asymptotic Distribution

To show the normality of  $V$ , we show that  $\frac{V_n - \mu_n}{\sigma_{V_n}} \sim N(0, 1)$  as  $n \rightarrow \infty$ .

We verify, as  $n \rightarrow \infty$ ,  $m_{2r} = \frac{(2r)!}{2^r r!}$  for every  $r$  and  $m_{2r-1} = 0$ .

# Asymptotic Distribution

## Corollary

*The distribution of number of  $V$  on a tiling of  $2$ -by- $n$  board is asymptotically normal.*

## Proof.

$$\frac{E[(V_n - \mu_n)^{2r}]}{E[(V_n - \mu_n)^2]^r} = \frac{\left(\frac{2}{5\sqrt{5}}\right)^r \frac{n^r}{r!} (2r)!}{\left(\frac{4n}{5\sqrt{5}}\right)^r} = \frac{(2r)!}{2^r r!},$$

and

$$\frac{E[(V_n - \mu_n)^{2r+1}]}{E[(V_n - \mu_n)^2]^{(r+1/2)}} = 0,$$



# Asymptotic Distribution



**Remark** The conjectures of the straight moment  $E[V^r]$  can be verified similarly with less calculation.

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# Statistics on an $m$ -by- $n$ board

The method we used for a 2-by- $n$  board can be generalized to an  $m$ -by- $n$  board for a fixed  $m$  and symbolic  $n$ .

# References

-  Doron Zeilberger, *Automated Derivation of Limiting Distributions Of Combinatorial Random Variables Whose Generating Functions are Rational*, The Personal Journal of Shalosh B. Ekhad and Doron Zeilberger, Dec. 24, 2016.
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